# Context-Free Parsing 

He He

New York University
October 27, 2021

## Logistics

- Homework 3 released
- Project proposal due next week (one page)
- What problem are you tackling and why is it important?
- What's your approach?
- How do you plan to evaluate it?


## Table of Contents

Context-free language

Probabilistic context-free grammars

Discriminative parsing

## Langauge is a set of strings

## Formal language:

- A set of strings consisting of words from an alphabet
- Well-formed according to a set of rules
- Studies the syntactical aspects of a language

Examples:

- Formulas (logic): $\left.\left(p_{1} \wedge p_{2}\right) \vee( \urcorner p_{3}\right)$
- Programming languages: int $\mathrm{a}, \mathrm{b}=0$;
- Sequences from the alphabet $\{a, b\}$ that ends with two a's

Questions:

- Formal language theory: How to describe languages (expressive power, recognizability etc.)
- Linguistics: Can we design formal languages that capture syntactic properties of natural language?


## Natural language syntax

Construct a formal language to represent the syntax of natural language

- Expressivity: how many syntactic phenomena can it cover?
- Computation: how fast can we parse a sentence?

Context-free grammars for natural language

- Captures nested structures which are common in natural language
[I told Mary that [John told Jane that [Ted told Tom a secret]]].
- Captures long-range dependencies
the burnt and badly-ground Italian coffee
these burnt and badly-ground Italian coffees
- Strikes a good balance between expressivity and computation


## Context-free language

Context-free languages (CFL) are generated by a context-free grammar $G=(\Sigma, N, R, S):$

- a finite alphabet $\Sigma$ of terminals (words)
- a finite set of non-terminals $N$ disjoint from $\Sigma$ (word groups)
- a set of production rules $R$ of the form $A \rightarrow \beta$, where $A \in N, \beta \in(\Sigma \cup N)^{*}$ (how to group words)
- a start symbol $S \in N$ (root of derivation)


## Example:

$$
\left.\begin{array}{l}
S \rightarrow S S \\
S \rightarrow(S) \\
S \rightarrow()
\end{array}\right]
$$


$(1)(c))$

## Phrase-structure grammar for English

Sentences are broken down into constituents.
A constituent works as a single unit in a sentence.

- Can be moved around or replaced without breaking grammaticality. (Abigail) and (her younger brother) (bought a fish).

Construct CFG for English

- Each word is a terminal, derived from its POS tag. DT $\rightarrow$ the NN $\rightarrow$ fish
- Each sentence is derived from the start symbol $S$.
- Each phrase type is a non-terminal.
- Each constituent is derived from a non-terminal.

Grammar design: choose the right set of non-terminals that produces different constituents.

## A toy example CFG

$$
\begin{aligned}
& N=\{\mathrm{S}, \mathrm{NP}, \mathrm{VP}, \mathrm{PP}, \mathrm{DT}, \mathrm{Vi}, \mathrm{Vt}, \mathrm{NN}, \mathrm{IN}\} \\
& S=\text { S } \\
& \Sigma=\{\text { sleeps, saw, man, woman, dog, telescope, the, with, in }\} \\
& R=
\end{aligned}
$$

Lexicon: rules that produce the terminals
(Example from Mike Collins' notes)

Parsing

$$
R=
$$

| S | $\rightarrow$ | NP | VP |
| :--- | :--- | :--- | :--- |
| VP | $\rightarrow$ | Vi |  |
| VP | $\rightarrow$ | Vt | NP |
| VP | $\rightarrow$ | VP | PP |
| NP | $\rightarrow$ | DT | NN |
| NP | $\rightarrow$ | NP | PP |
| PP | $\rightarrow$ | IN | NP |


| Vi | $\rightarrow$ | sleeps |
| :--- | :--- | :--- |
| Vt | $\rightarrow$ | saw |
| NN | $\rightarrow$ | man |
| MN | $\rightarrow$ | woman |
| ND | $\rightarrow$ | telescope |
| KN | $\rightarrow$ | dog |
| OT | $\rightarrow$ | the |
| IN | $\rightarrow$ | with |
| IN | $\rightarrow$ | in |

Can we derive the sentence "the man sleeps"?
(1)

(2) [s [up lot the ] [uv man] ] [up [vi sleeps $)$ ] $]$
(3) $s_{1}=S \rightarrow N P \vee P$
$\rightarrow$ DT NN vP
$\rightarrow$ The NN $r$ P
$i$

## Ambiguity

Can a sentence have multiple parse trees?


Exercise: find parse trees for
"She announced a program to promote safety in trucks and vans".

## Table of Contents

## Context-free language

Probabilistic context-free grammars

Discriminative parsing

## PCFG

Notation: let $\mathcal{T}_{G}$ be the set of all possible left-most parse trees under the grammar $G$.
Goal: define a probability distribution $p(t)$ over parse trees $t \in \mathcal{T}_{G}$
Parsing: pick the most likely parse tree for a sentence $s$

$$
\underset{t \in \mathcal{T}_{G}(s)}{\arg \max } p(t)
$$

Three questions:

- Modeling: how to define $p(t)$ for trees?
- Learning: how to estimate parameters of the distribution $p(t)$ ?
- Inference: how to find the most likely tree efficiently?


## Modeling

Generate parse trees: iteratively sample a production rule to expand a non-terminal

$$
R=
$$

| S | $\rightarrow$ | NP | VP |
| :--- | :--- | :--- | :--- |
| VP | $\rightarrow$ | Vi |  |
| VP | $\rightarrow$ | Vt | NP |
| VP | $\rightarrow$ | VP | PP |
| NP | $\rightarrow$ | DT | NN |
| NP | $\rightarrow$ | NP | PP |
| PP | $\rightarrow$ | IN | NP |



## PCFG

## A PCFG consists of

- A CFG $G=(\Sigma, N, R, S)$
- Probabilities of production rules $q(\alpha \rightarrow \beta)$ for each $\alpha \rightarrow \beta \in R$ such that

$$
\sum_{\beta: X \rightarrow \beta \in R} q(X \rightarrow \beta)=1 \quad \forall X \in N
$$

$R, q=$

| S | $\rightarrow$ | NP | VP | 1.0 |
| :--- | :--- | :--- | :--- | :--- |
| VP | $\rightarrow$ | Vi |  | 0.3 |
| VP | $\rightarrow$ | Vt | NP | 0.5 |
| VP | $\rightarrow$ | VP | PP | 0.2 |
| NP | $\rightarrow$ | DT | NN | 0.8 |
| NP | $\rightarrow$ | NP | PP | 0.2 |
| PP | $\rightarrow$ | IN | NP | 1.0 |

$\left.\begin{array}{|lll|l|}\hline \mathrm{Vi} & \rightarrow & \text { sleeps } & 1.0 \\ \mathrm{Vt} & \rightarrow & \text { saw } & 1.0 \\ \hline \mathrm{NN} & \rightarrow & \text { man } & 0.1 \\ \mathrm{NN} & \rightarrow & \text { woman } & 0.1 \\ \mathrm{NN} & \rightarrow & \text { telescope } & 0.3 \\ \mathrm{NN} & \rightarrow & \text { dog } & 0.5 \\ \hline \mathrm{DT} & \rightarrow & \text { the } & 1.0 \\ \hline \mathrm{IN} & \rightarrow & \text { with } & 0.6 \\ \mathrm{IN} & \rightarrow & \text { in } & 0.4 \\ \hline\end{array}\right]=1$

From HMM to PCFG

## Probabilities of parse trees

Given a parse tree $t$ consisting of rules $\alpha_{1} \rightarrow \beta_{1}, \ldots, \alpha_{n} \rightarrow \beta_{n}$, its probabilities under the PCFG is

$$
p(t)=\prod_{i=1}^{n} q\left(\alpha_{i} \rightarrow \beta_{i}\right)
$$

## Example:



## Learning

Training data: treebanks


Given a set of trees (production rules), we can estimate rule probabilities by MLE.

$$
q(\alpha \rightarrow \beta)=\frac{\operatorname{count}(\alpha \rightarrow \beta)}{\sum_{\beta^{\prime}: \alpha \rightarrow \beta^{\prime} \in R} \operatorname{count}\left(\alpha \rightarrow \beta^{\prime}\right)} \quad \begin{aligned}
& \text { S } \rightarrow \text { NP VP } \\
& \mathrm{S} \rightarrow *
\end{aligned}
$$

- Similar to estimate word probabilities $(\rightarrow \beta)$ given the document class $(\alpha)$.


## Parsing

Input: sentences, (P)CFG
Output: derivations / parse trees (with scores/probabilities)
Total number of parse trees for a sentence?
Consider a minimal CFG:

$$
\frac{X \rightarrow X X}{X \rightarrow \text { aardvark|abacus }|\ldots| \text { zyther }}
$$



Given a string, \# of parse trees $=\#$ of strings with balanced brackets

$$
\left(\left(w_{1} w_{2}\right)\left(w_{3} w_{4}\right)\right),\left(\left(\left(w_{1} w_{2}\right) w_{3}\right) w_{4}\right), \ldots
$$

\# of strings with $n$ pairs of brackets:

$$
\text { Catalan number } C_{n}=\frac{1}{n+1}\binom{2 n}{n}
$$

## Chomsky normal form (CNF)

A CFG is in Chomsky normal form if every production rule takes one of the following forms:

- Binary non-terminal production: $A \rightarrow B C$ where $A, B, C \in N$.
- Unary terminal production: $A \rightarrow a$ where $A \in N, a \in \Sigma$.

Grammars in CNF produces binary parse trees.
Binarize a production rule: VP $\rightarrow$ VBD NP PP

$$
\begin{aligned}
& \mathrm{VP} \rightarrow \text { VBD@VP-VBD } \\
& \text { @VP-VBD } \rightarrow \text { NP PP }
\end{aligned}
$$

We assume the grammar are in CNF.

## Dynamic programming on the tree

$$
p(t)=\underbrace{q(A \rightarrow B C)}_{\text {top rule }} \times \underbrace{p\left(t_{B}\right)}_{\text {left child }} \times \underbrace{p\left(t_{C}\right)}_{\text {right child }}
$$



What are the variables when constructing a tree rooted at $A$ spanning $x_{i}, \ldots, x_{j}$ ?

- The production rule $A \rightarrow B C$
- The splitting point $s$ : $B$ spans $x_{i}, \ldots, x_{s}$ and $C$ spans $x_{s+1}, \ldots, x_{j}$


## The CYK algorithm

Notation: $\mathcal{T}(i, j, X)$ is the set of trees with root node $X$ spanning $x_{i}, \ldots, x_{j}$
Subproblem:

$$
\pi(i, j, X)=\max _{t \in \mathcal{T}(i, j, X)} p(t)
$$

Base case:

$$
\pi(i, i, X)= \begin{cases}q\left(X \rightarrow x_{i}\right) & \text { if } X \rightarrow x_{i} \in R \\ 0 & \text { otherwise }\end{cases}
$$

Recursion:

$$
\pi(i, j, X)=\max _{\substack{Y, Z \in N \\ s \in\{i, \ldots, j-1\}}} q(X \rightarrow Y Z) \times \pi(i, s, Y) \times \pi(s+1, j, Z)
$$

Use backtracking to find the argmax tree.

Bottom-up parsing

## Variants of CYK

Argmax: find the most likely tree (analogous to Viterbi).

$$
\pi(i, j, X)=\max \underset{s \in\{i, \ldots, j-1\}}{Y, Z \in N} q(X \rightarrow Y Z) \times \pi(i, s, Y) \times \pi(s+1, j, Z)
$$

Recognition: does the string belong to the language?

$$
\pi(i, j, X)=\vee \underset{s \in\{i, \ldots, j-1\}}{Y, Z \in N} \mathbb{I}[X \rightarrow Y Z \in R] \wedge \pi(i, s, Y) \wedge \pi(s+1, j, Z)
$$

Marginalization: what's the probability of the string being generated from the grammar? (the inside algorithm) $\quad \sum_{z} p(x, z)$

$$
\pi(i, j, X)=\sum_{\substack{ \\s \in\{i, \ldots, j-1\}}} q(X \rightarrow Y Z) \times \pi(i, s, Y) \times \pi(s+1, j, Z)
$$

## Summary

|  | NB | HMM | PCFG |
| :--- | :--- | :--- | :--- |
| output structure | category | sequence | tree |
| learning |  | MLE |  |
| decoding | bruteforce | Viterbi <br> $p\left(y_{i} \mid x\right)$, <br> marginalization | $p\left(y_{i}, y_{i-1} \mid x\right)$ |$\quad p(i, j, N \mid x)$

## Table of Contents

Context-free language

Probabilistic context-free grammars

Discriminative parsing

## CRF for trees

Input: sequence of words $x=\left(x_{1}, \ldots, x_{n}\right)$
Output: parse tree $y \in \mathcal{T}(x)$
Model: decompose by production rules

$$
p(y \mid x ; \theta) \propto \prod_{(r, s)} \psi(r, s \mid x ; \theta)
$$

- $r$ : production rule
- $s$ : start, split, end indices of the rule $r$



## CRF parsing

Potential functions:

$$
\begin{aligned}
\psi(r, s \mid x ; \theta) & =\exp (\theta \cdot \phi(r, s, x)) \\
\prod_{(r, s) \in \mathcal{T}(x)} \psi(r, s \mid x ; \theta) & =\exp \left(\sum_{(r, s) \in \mathcal{T}(x)} \theta \cdot \phi(r, s, x)\right)
\end{aligned}
$$

Learning: MLE

1. Compute the partition function by the inside algorithm
2. Call autograd to compute the gradient (backpropagation)

Inference: CYK

## Limitations of PCFG



Limited lexical information

## Lexicalized PCFG

Attach the "head" (most important child in a rule) of the span to each non-terminal


## Features

Easy to incorporate lexical information in features!

$$
\text { local score }=\theta \cdot \phi(\mathrm{VP} \rightarrow \text { VBD NP, }(5,6,8), \ldots \text { averted financial disaster... })
$$



Figure: Less grammar, more features. [Hall+ 14]

## Neural CRF parser


reflected the flip side of the Stoltzman personality


Figure: Neural CRF Parsing. [Durrett+ 15]

- $f_{w}$ : lexical features
- $f_{o}$ : rule features
- $h^{T} W f_{o}$ : interaction between lexical and rule features


## Evaluation

$$
\begin{aligned}
\text { recall } & =\frac{\# \text { correct constituents }}{\# \text { total constituents in gold trees }} \\
\text { precision } & =\frac{\# \text { correct constituents }}{\# \text { total constituents in predicted trees }} \\
\text { F1 } & =\frac{2 \times \text { precision } \times \text { recall }}{\text { precision }+ \text { recall }}
\end{aligned}
$$

- Constituent: $(i, j, X)$
- Labeled F1: the non-terminal node label must be correct
- Unlabeled F1: just consider the tree structure


## Example



